INTRODUCTION

The 24XXX series serial EEPROMs from Microchip Technology support a bidirectional, 2-wire bus and data transmission protocol. The bus is controlled by the microcontroller (master), which generates the Serial Clock (SCL), controls the bus access and generates the Start and Stop conditions, while the 24XXX serial EEPROM works as slave. The 24XXX serial EEPROMs are \( \text{I}^2\text{C}\)™ compatible and have maximum clock frequencies ranging from 100 kHz to 1 MHz.

The main features of the 24XXX serial EEPROMs are:

- 2-wire serial interface bus, \( \text{I}^2\text{C}\) compatible
- EEPROM densities from 128 bits to 512 Kbits
- Bus speed from 100 kHz to 1 MHz
- Voltage range from 1.7V to 5.5V
- Low power operation
- Temperature range from -40°C to +125°C
- Over 1,000,000 erase/write cycles
- Up to 8 devices may be connected to same bus

Figure 1 is the hardware schematic depicting the interface between the Microchip 24XXX series of \( \text{I}^2\text{C}\) serial EEPROMs and NXP’s P89LPC952 8051-based MCU. The schematic shows the connections necessary between the MCU and the serial EEPROM as tested, as well as the required pull-up resistors on the clock line (SCL) and data line (SDA). Not illustrated in this application note are the write-protect feature and the cascading of multiple devices; thus, the WP pin and address pins A0, A1 and A2 are tied to Vss (ground). The test software was written assuming these connections.

Note: A decoupling capacitor (typically 0.1 \( \mu\text{F} \)) should be used to filter noise on Vcc.
FIRMWARE DESCRIPTION

Main Function

The purpose of the firmware is to show how to generate specific I2C bus transactions using the bidirectional SDA pin on the microcontroller. The focus is to provide the designer with a strong understanding of communication with the 24XXX series serial EEPROMs, thus allowing for more complex programs to be written in the future. The firmware was written in C for NXP’s P89LPC952 MCU using the Keil™ µVision3® IDE and was developed on the Keil MCB950 evaluation board.

The main code demonstrates two different methods of accessing the I2C serial EEPROM: byte access and page access. The byte method accesses single bytes, where every data byte is preceded by three bytes of address: device address, MSB address, and LSB address. In the page access method, the MCU sends the address of the first byte, and the I2C serial EEPROM internally increments the address pointer for the next data byte.

The code was tested using the 24XX512 serial EEPROM. The EEPROM features 64K x 8 (512 Kbit) of memory and a page write capability of up to 128 bytes of data. Oscilloscope screen shots are shown in this application note. All timings are based on the internal RC oscillator of the MCU (7.373 MHz). If a faster clock is used, the code must be modified to generate the correct delays.

The bus speed in these examples is ~75 kHz.

I2C Functions

When an MCU accesses an I2C serial EEPROM, it is always the master on the I2C bus and the I2C serial EEPROM is the slave. The MCU controls all operations on the bus. Each operation is started by the MCU through a Start condition, followed by a control byte. The control byte consists of the control code (first 4 bits), the device address (next 3 bits), and the read/write (R/W) bit. The control code is always the same for the serial EEPROM being accessed, while the device address can range from '000' to '111', allowing up to eight different devices on the same bus. The R/W bit tells the serial EEPROM which operation to perform.

To access an I2C serial EEPROM at the start, the MCU writes the device address and the byte address to the I2C serial EEPROM; thus, each access cycle starts with a Write condition. For read operations, after the above sequence, the MCU switches from Transmitter mode to Receiver mode and the serial EEPROM from Receiver to Transmitter mode through a Restart condition.

BYTE WRITE OPERATION

Figure 2 depicts the necessary components that comprise the byte write operation. Each MCU’s action is acknowledged (ACK) by the I2C serial EEPROM on the 9th bit of the clock by pulling down the SDA data line; consequently, every byte transfer lasts for 9 clock transitions.

FIGURE 2: BYTE WRITE

<table>
<thead>
<tr>
<th>Bus Activity</th>
<th>MCU</th>
<th>SDA Line</th>
<th>Bus Activity</th>
</tr>
</thead>
<tbody>
<tr>
<td>S</td>
<td>T</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>1</td>
<td>0</td>
<td>1</td>
</tr>
<tr>
<td>T</td>
<td>A</td>
<td>2</td>
<td>1</td>
</tr>
<tr>
<td>R</td>
<td>P</td>
<td>STOP</td>
<td></td>
</tr>
</tbody>
</table>

Control Byte/Device Address | MSB Address | LSB Address | Data Byte
BYTE READ OPERATION

Figure 3 depicts the necessary components that comprise the byte read operation. The second Start condition instructs the I^2C serial EEPROM to place data on the I^2C bus.

The SDA line must remain stable while the SCL clock line is high. Any change of the SDA line while the SCL line is high is interpreted by the I^2C serial EEPROM as a Start or Stop condition.

FIGURE 3: BYTE READ

PAGE WRITE OPERATION

Figure 4 depicts the necessary components that comprise the page write operation. The only difference between the page write operation and the byte write operation (Figure 2) is that the MCU, instead of sending 1 byte, sends 'n' bytes of data, up to the maximum page size of the I^2C serial EEPROM.

FIGURE 4: PAGE WRITE

SEQUENTIAL READ OPERATION

Figure 5 depicts the necessary components that comprise the sequential read operation. The last read byte is not acknowledged (NACK) by the MCU. This terminates the sequential read operation.

FIGURE 5: SEQUENTIAL READ
**INITIALIZATION**

Initialization consists of initializing the hardware peripheral within the MCU. It performs the following actions:

- I2EN = 1; //enable I2C part
- CRSEL = 0; //select internal baud generator
- STA = STO = SI = AA = 0; //clear Start, Stop, serial flag, ACK bit
- Sets the bus speed at ~75 kHz: I2SCLL = I2SCLH = 25

The bit frequency \( f_{bit} \) will be:

\[
f_{bit} = \frac{f_{PCLK}}{2 \times (I2SCLH + I2SCLL)}
\]

**EQUATION 1: BIT RATE**

**START DATA TRANSFER**

The MCU generates a Start condition on the I2C bus to initiate data transfer.

Figure 6 shows a typical scope plot from the beginning of a write operation. Any memory access begins with a Start condition. This is followed by the I2C serial EEPROM (slave) address (A0h). Because the R/W bit is set to ‘0’, the next operation of the bus is a write.
STOP DATA TRANSFER

The MCU generates a Stop condition on the I²C bus to stop data transfer.

Figure 7 shows a typical scope plot of a byte write operation followed by a Stop condition. Every operation on the I²C bus ends with a Stop condition.

FIGURE 7: BYTE WRITE AND STOP CONDITION
REPEATED START
The Repeated Start condition switches the MCU from Transmitter to Receiver mode and the serial EEPROM from Receiver to Transmitter. The bit is followed by the write of the \( \text{I}^\text{2C} \) device address and R/W bit sequence. The sequence is presented in Figure 8.

WRITE \( \text{I}^\text{2C} \) SERIAL EEPROM ADDRESS AND READ CONDITION
Following the Repeated Start condition, the MCU writes the \( \text{I}^\text{2C} \) serial EEPROM address (slave address) with the R/W bit set correctly.

For read sequences, the \( \text{I}^\text{2C} \) serial EEPROM address is A1h because the R/W bit is ‘1’, indicating that the MCU waits to read bytes from the \( \text{I}^\text{2C} \) serial EEPROM.

The sequence is presented in Figure 8.
WRITE DEVICE ADDRESS AND ADDRESS BYTES

After the Start condition, the MCU writes three bytes, consisting of the device address and the MSB and LSB addresses. Because the device address is followed by the write of the two address bytes, the device address has the R/W bit set to '0'.

The scope plot in Figure 9 depicts the MSB address byte (00h), the LSB address byte (60h), followed by the first written byte (through the byte access method) of the string (43h). Because the byte access method accesses single bytes, the data byte is followed by a Stop bit.

The short spikes observed in the scope plot represent a short period of time during which both the MCU and the I²C serial EEPROM release the data bus.

FIGURE 9: WRITE MSB AND LSB ADDRESS BYTES
WRITE 8 BITS OF DATA

The write data function is used in both byte write and page write operations. The structure of the byte write operation is shown in Figure 2. The structure of the page write operation is shown in Figure 4. The only difference between the two operations is that in the page write operation the MCU sends ‘n’ bytes of data instead of only one.

As a rule, the function shifts from a parallel format to the serial \( I^2C \) format. The bus speed is in accordance with the initialization routine and is \( \sim 75 \text{ kHz} \).

The MCU sets the data line on the falling edge of the clock, and the \( I^2C \) serial EEPROM latches this in on the rising edge of the clock.

In Figure 6 a spike labeled “bus release” can be seen between the 9th clock pulse and the next clock pulse. The spike is the sign that both devices – the MCU and the \( I^2C \) serial EEPROM – released the open-drain SDA line in order to be able to continue the communication.

READ 8 BITS OF DATA

The read data function is used in both byte read and sequential read operations. The structure of the byte read operation is shown in Figure 3. The structure of the sequential read operation is shown in Figure 5.

In the byte read operation, the 8-bit read data is situated at the end of the command and is not acknowledged by the MCU through a NACK on the 9th clock pulse of the read data byte, which is shown in Figure 10.

In the case of a sequential read operation, the last read byte will be not acknowledged and the other \(<n-1>\) read bytes will be acknowledged by the MCU.

FIGURE 10:  BYTE READ, NACK AND STOP CONDITION
WRITE DATA BYTE

The structure of this byte write operation is shown in Figure 2. Basically the function consists of a sequence of function calls. The start data transfer sequence is described in detail in the section entitled “Start Data Transfer” and in Figure 6. The stop data transfer sequence is described in detail in the section entitled “Stop Data Transfer” and in Figure 7. The write MSB and LSB address byte sequence is described in the section entitled “Write Device Address and Address Bytes” and in Figure 9.

READ DATA BYTE

The read data byte function reads a data byte from the I2C serial EEPROM. The structure of the operation is shown in Figure 3.

After the first Start condition, the MCU sends the device address, the MSB address byte, then the LSB address byte to the I2C serial EEPROM.

Once the entire memory address has been sent to the serial EEPROM, a new Start condition is generated to switch the MCU from Transmitter to Receiver mode and the serial EEPROM from Receiver to Transmitter mode. Before the read, the MCU must send a new device address for a read.

The MCU must generate the necessary NACK or ACK conditions to terminate or continue the bus operation.

WRITE A STRING (PAGE WRITE)

In this application note, the length of the string is 16 bytes and the physical page size for the 24XX512 is 128 bytes. The length of the written string must be shorter than the physical page size. If the page write operation overwrites the physical page boundary, the internal address counter rolls over and overwrites the first bytes of the current page.

The structure of the page write operation is shown in Figure 4.

The sequence must send the address of the first byte to be written. The address is automatically incremented at every byte write by the internal logic of the I2C serial EEPROM. Each received byte is acknowledged by the EEPROM.

The scope plot in Figure 11 depicts the write of the first three consecutive bytes.
FIGURE 11: PAGE WRITE (FIRST 3 BYTES)

Bus Activity

MCU

SDA Line

Bus Activity

Data Byte 0

Data Byte 127
READ A STRING (SEQUENTIAL READ)

In contrast to the page write operation described in the previous paragraph, there is no maximum length for a sequential read. After 64 Kbytes have been read, the internal address counter rolls over to the beginning of the array.

To indicate the end of the read, the MCU sends a NACK before the Stop condition. All other previously read bytes are acknowledged by the MCU.

The structure of the sequential read operation is shown in Figure 5. Figure 12 shows a typical scope plot depicting this operation.

FIGURE 12: SEQUENTIAL READ (FIRST BYTES)
BYTE WRITE VERSUS PAGE WRITE

At first glance, the page write method appears superior to the byte write method: it’s simpler and faster. However, a careful analysis shows that the byte write method has a major advantage over page write owing to the roll-over phenomenon (see Note).

As a consequence of the roll-over phenomenon, applications that write long strings to the I²C serial EEPROM risk overlapping the page boundary in the middle of a string. In such instances, the firmware should use byte write to avoid this condition. The disadvantage of doing this is the slower speed involved in writing the entire string: every byte write cycle time is approximately 5 ms.

The following summarizes the differences between the byte write and page write methods.

**Byte Write**
- Is slower – It needs a 5 ms write cycle time for each byte.
- Is more general – It may write a string of any length.

**Page Write**
- Is faster – It needs only one write cycle time for the whole page.
- Care must be taken to observe page boundaries during page writes.

**CONCLUSION**

This application note offers designers a set of firmware routines to access I²C serial EEPROMs using a hardware peripheral. The code demonstrates byte and page operations. All routines were written in C for an 8051-based MCU.

The code was developed on the Keil MCB950 evaluation board using the schematic shown in Figure 1. It was tested using the NXP P89LPC952 MCU and debugged using the Keil µVision3 IDE.

Note: Page write operations are limited to writing bytes within a single physical page, regardless of the number of bytes actually being written. Physical page boundaries start at addresses that are integer multiples of the page buffer size (or page size), and they end at addresses that are integer multiples of [page size-1]. If a Page Write command attempts to write across a physical page boundary, the result is that the data wraps around to the beginning of the current page (overwriting data previously stored there) instead of being written to the next page as might be expected. It is therefore necessary for the application software to prevent page write operations that would attempt to cross a page boundary.
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